

Exact Simulation Infinite Server Queue

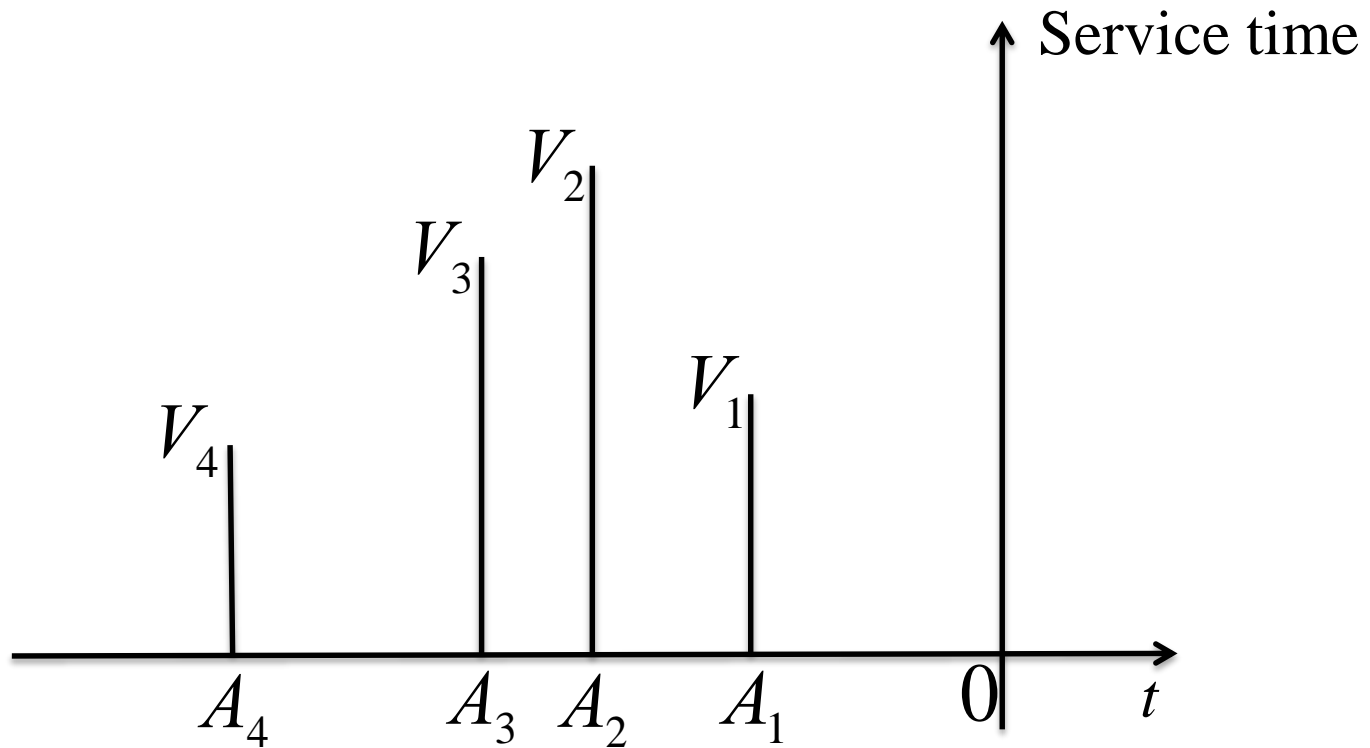
Jose Blanchet

Columbia University

(Joint work with Jing Dong.)

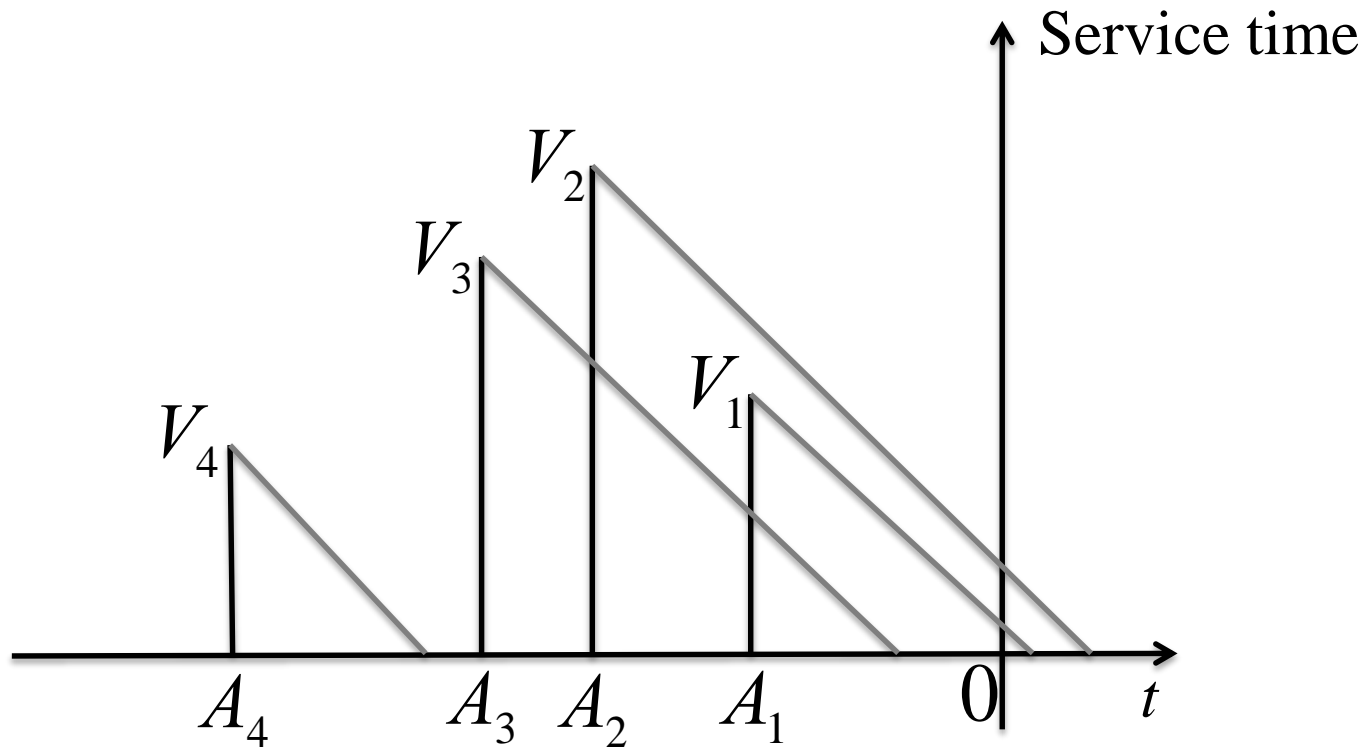
- I. Bias Reduction Techniques for Stochastic Networks
 - i. Point processes on stable unbounded regions
 - ii. Stationary infinite server queue
 - iii. Stationary loss systems

Infinite server queue



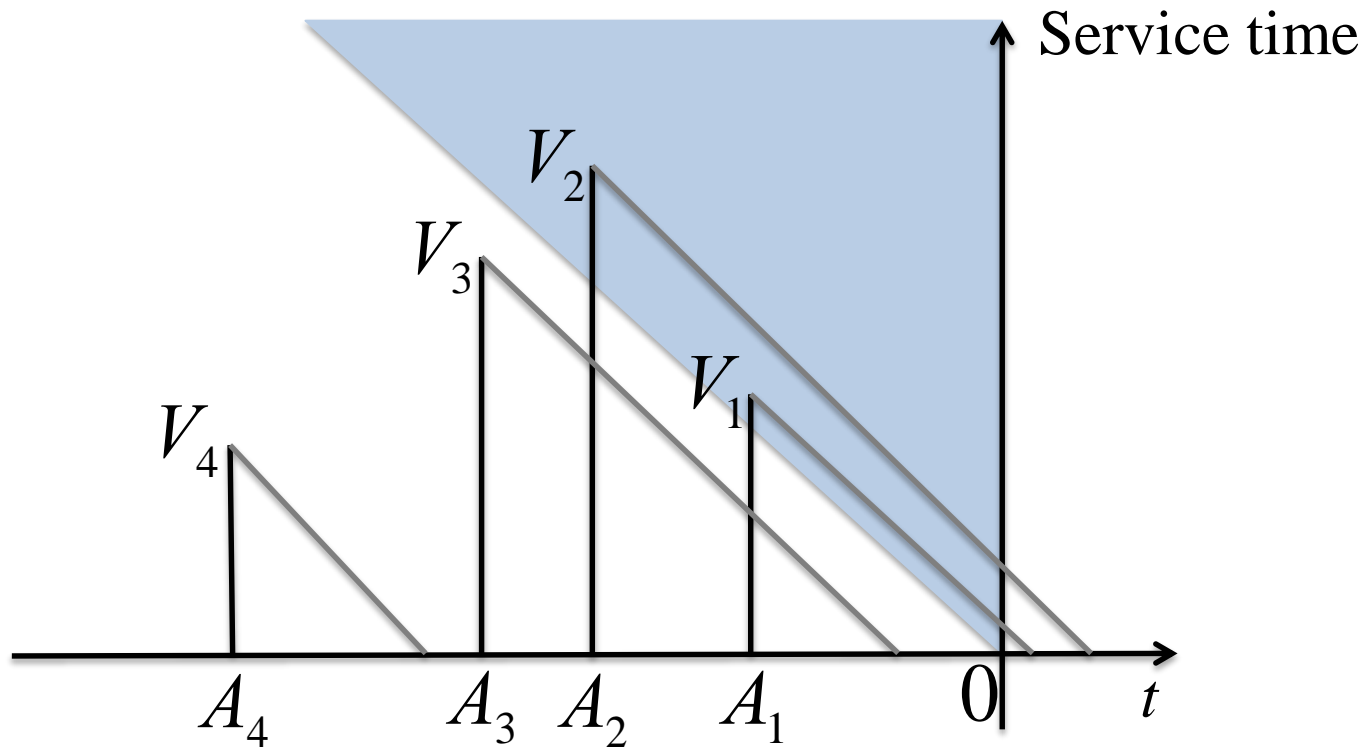
$$|A_n| = X_1 + X_2 + \dots + X_n$$

Infinite server queue



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Infinite server queue

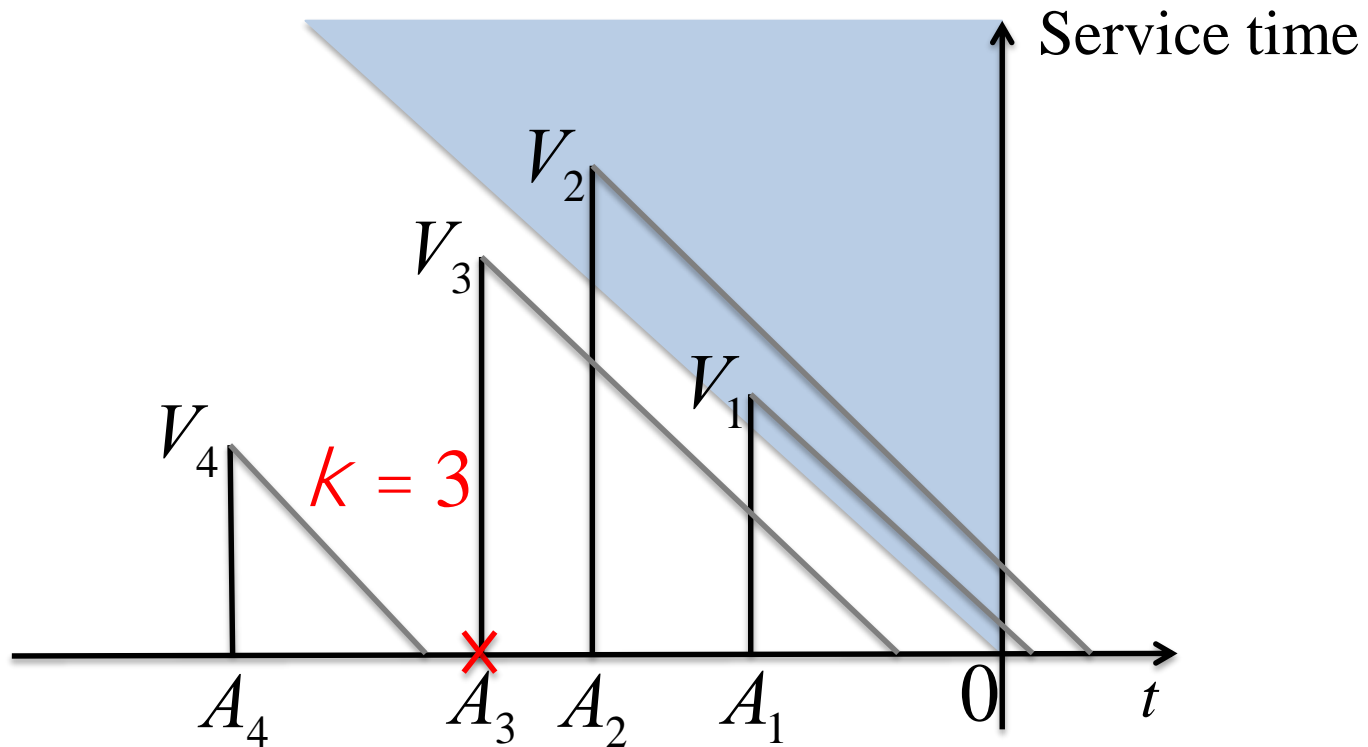


Record Breakers

Look into the future

- Define a sequence of record breakers
- Asking future yes/no questions: are there more record breakers or not
- If yes, find the next record breaker

Infinite server queue



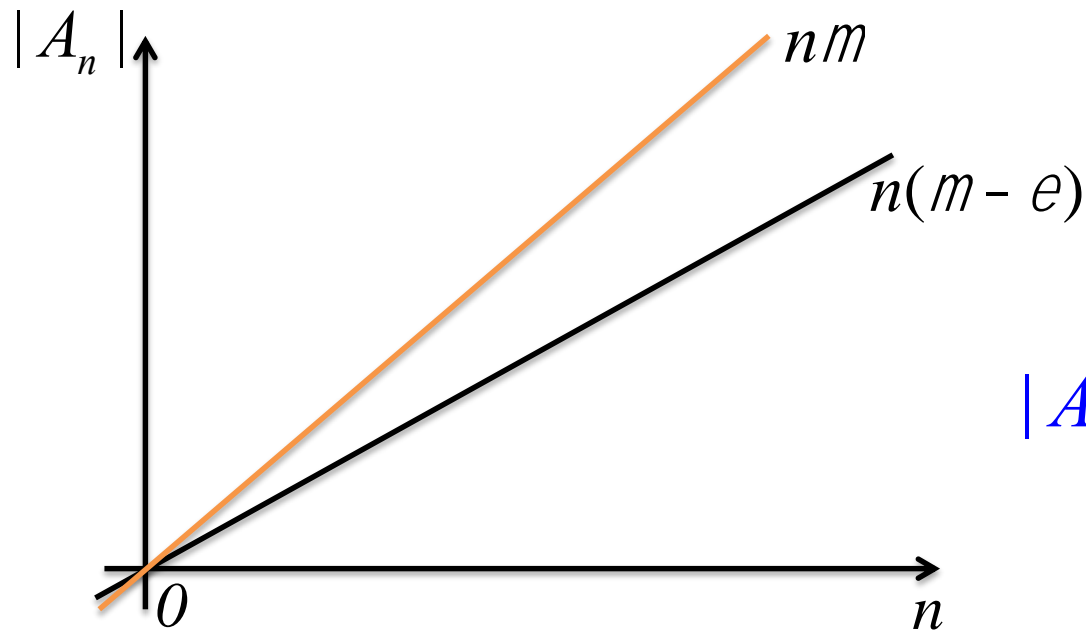
$$k : V_n \in |A_n| \text{ for all } n \geq k$$

Not a stopping time!

Infinite server queue

$$E[X_n] = m \text{ and fix } e \in (0, m)$$

$$k(A) : |A_n| \leq n(m - e) \text{ for all } n \in \mathbb{N}$$

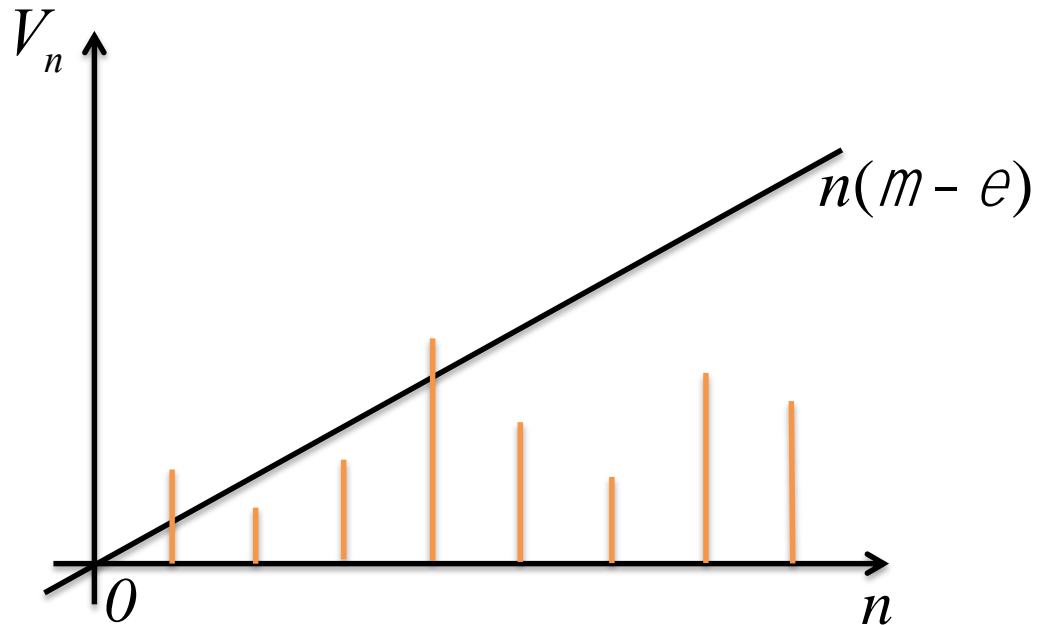


$$|A_n| = \sum_{i=1}^n X_i$$

Infinite server queue

$$E[X_n] = m \text{ and fix } e \in (0, m)$$

$$k(V) : V_n \leq n(m - e) \text{ for all } n \in \mathbb{N}$$



Infinite server queue

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$$k(A) : |A_n| \leq n(m - e) \text{ for all } n \geq k(A)$$

$$k(V) : V_n \leq n(m - e) \text{ for all } n \geq k(V)$$

$$k = \max\{k(V), k(A)\}$$

Infinite server queue: service time process

$$k(V) : V_n \leq n(m - e) \text{ for all } n \geq k(V)$$

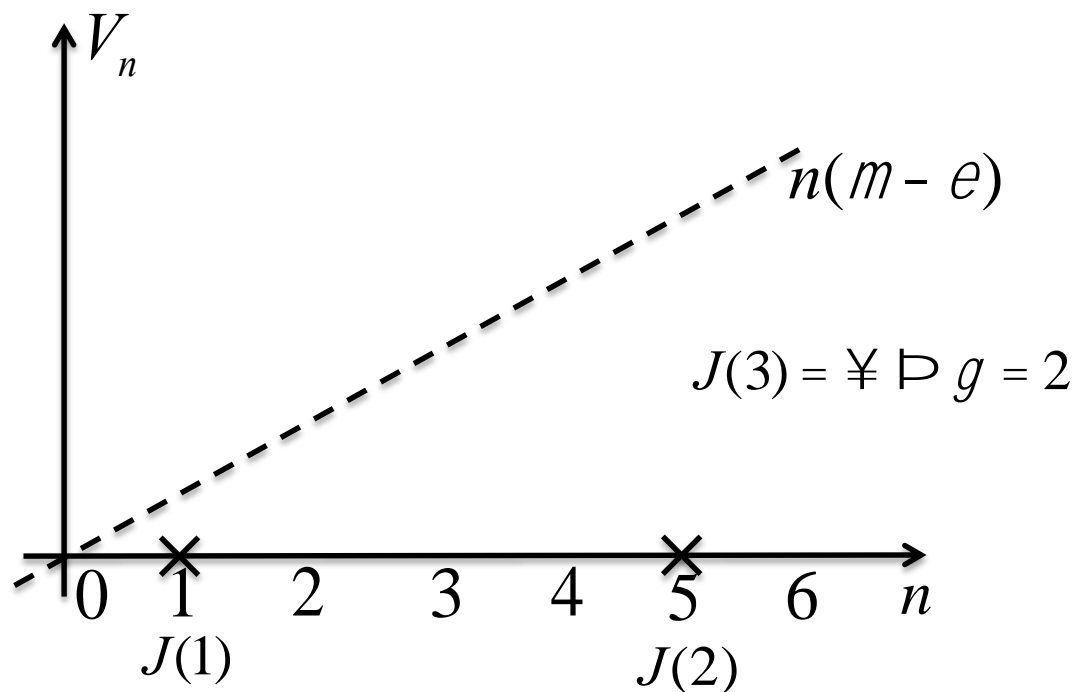
Record breaker: $\{n : V_n > n(m - e)\}$

$$J(0) := 0$$

$J(l)$: l -th record breaker

$J(g)$: last record breaker

$$\mathbb{P} k(V) = J(g) + 1$$



Infinite server queue: service time process

$k(V) : V_n \leq n(m - e)$ for all $n \geq k(V)$

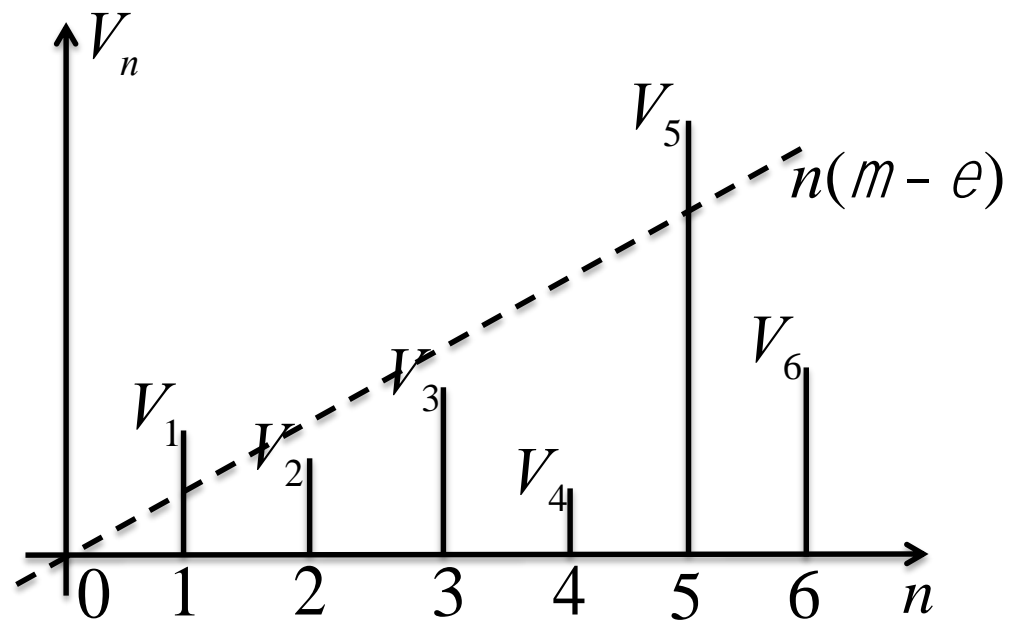
Record breaker: $\{n : V_n > n(m - e)\}$

$J(0) := 0$

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$$\vdash k(V) = J(g) + 1$$



Infinite server queue: service time process

Let $p(n) = P(V_i > n(m - e))$

Then $P(J(1) = \infty) = \prod_{n=1}^{\infty} (1 - p(n)) > 0$

Upper bound $P(J(1) = \infty) \leq \prod_{n=1}^h (1 - p(n)) := u(h)$

Lower bound $P(J(1) = \infty) \geq \prod_{n=1}^h (1 - p(n))g(h) := l(h)$

and $u(h) - u(h - 1) = p(h) \prod_{n=1}^{h-1} (1 - p(n)) = P(J(1) = h)$

Negative drifted random walk

$$S_n = X_1 + X_2 + \dots + X_n \text{ with } E[X_1] < 0$$

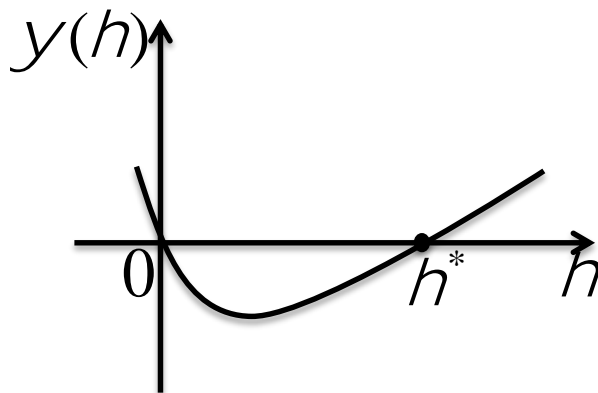
$$T_a = \inf\{n \geq 0 : S_n \geq a\}$$

$$P(T_a = \infty) > 0$$

Negative drifted random walk

$$f_h(y) = \exp(hy - y(h))f(y) \text{ where } y(h) := \log E[\exp(hY_i)]$$

$$h^* : y(h^*) := \log E[\exp(h^*Y_i)] = 0$$



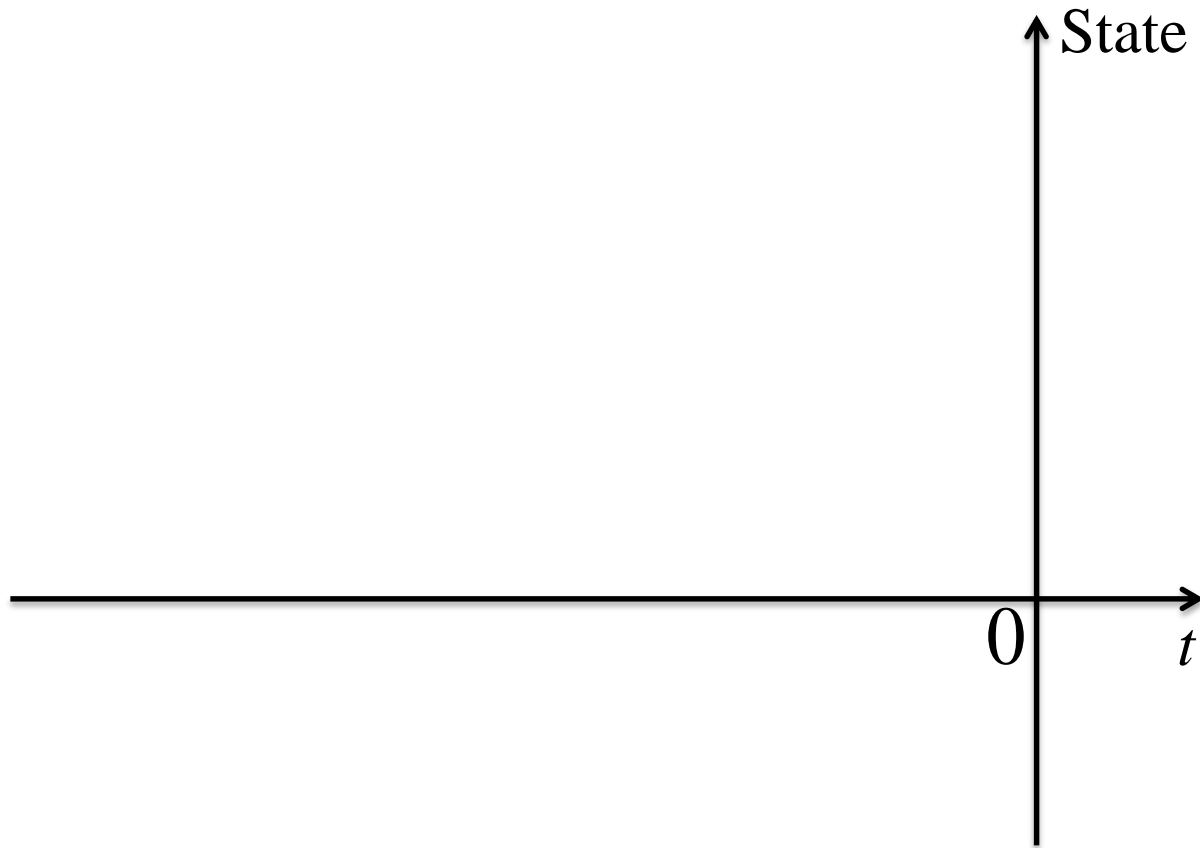
$$f_{h^*}(y) = f(y)\exp(h^*y)$$

$$P(T_a < \infty) = E_{h^*}[\exp(-h^*S_{T_a})1\{T_a < \infty\}] = E_{h^*}[\exp(-h^*S_{T_a})]$$

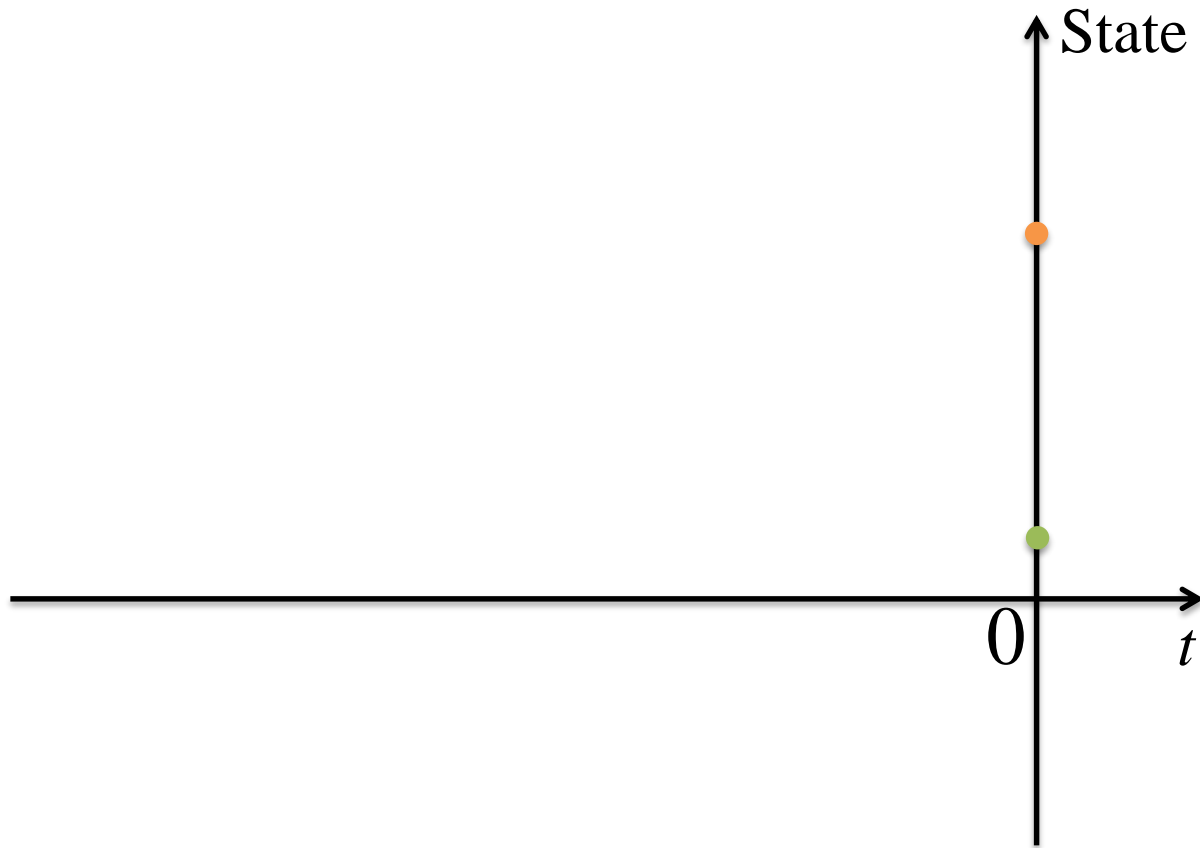
$$U \notin \exp(-h^*S_{T_a})$$

Dominated Coupling from the Past

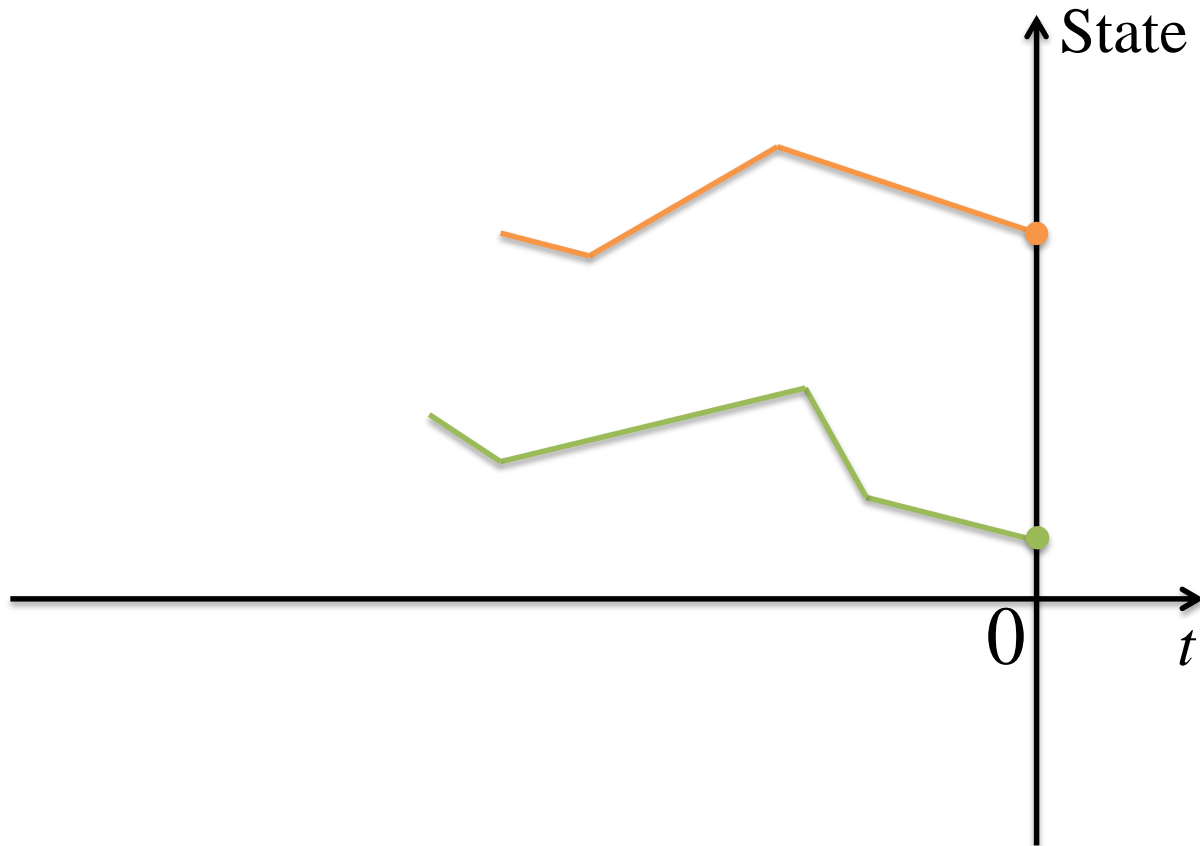
Dominated Coupling From The Past



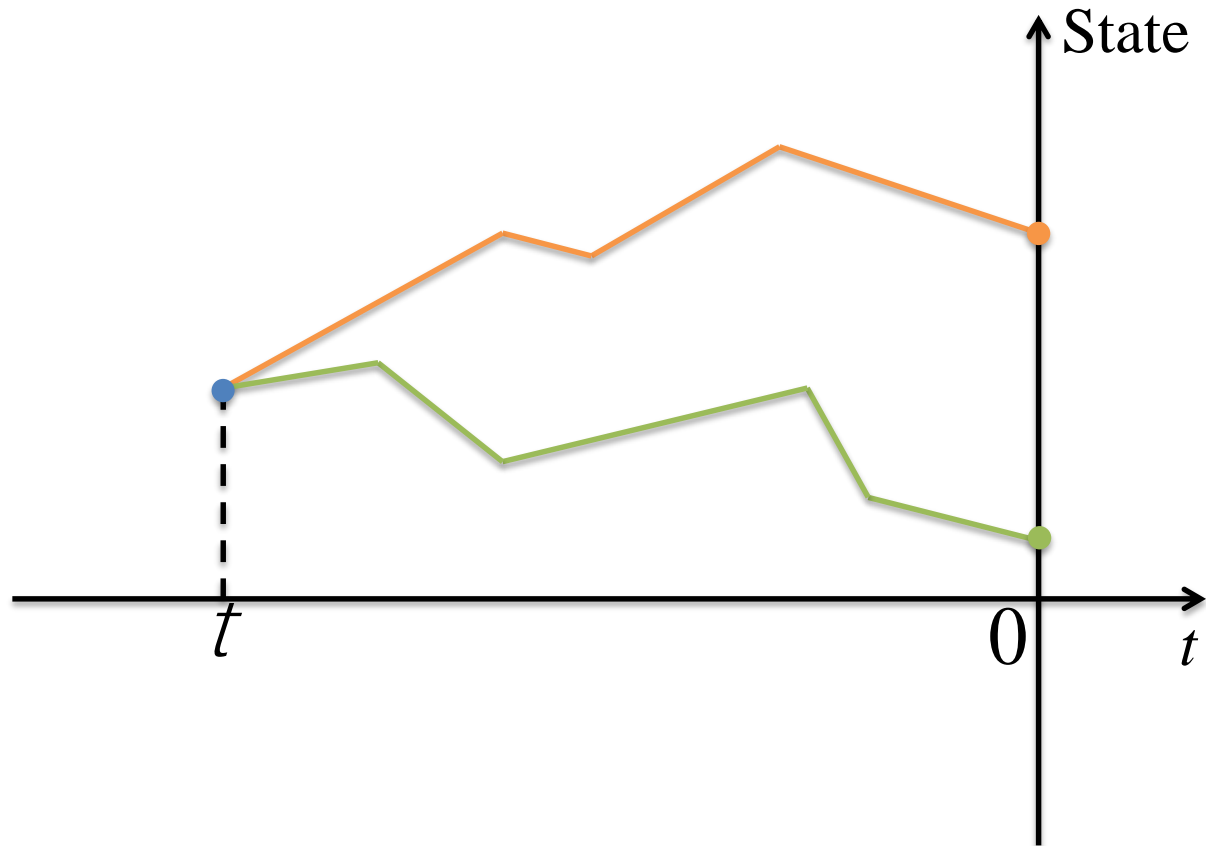
Dominated Coupling From The Past



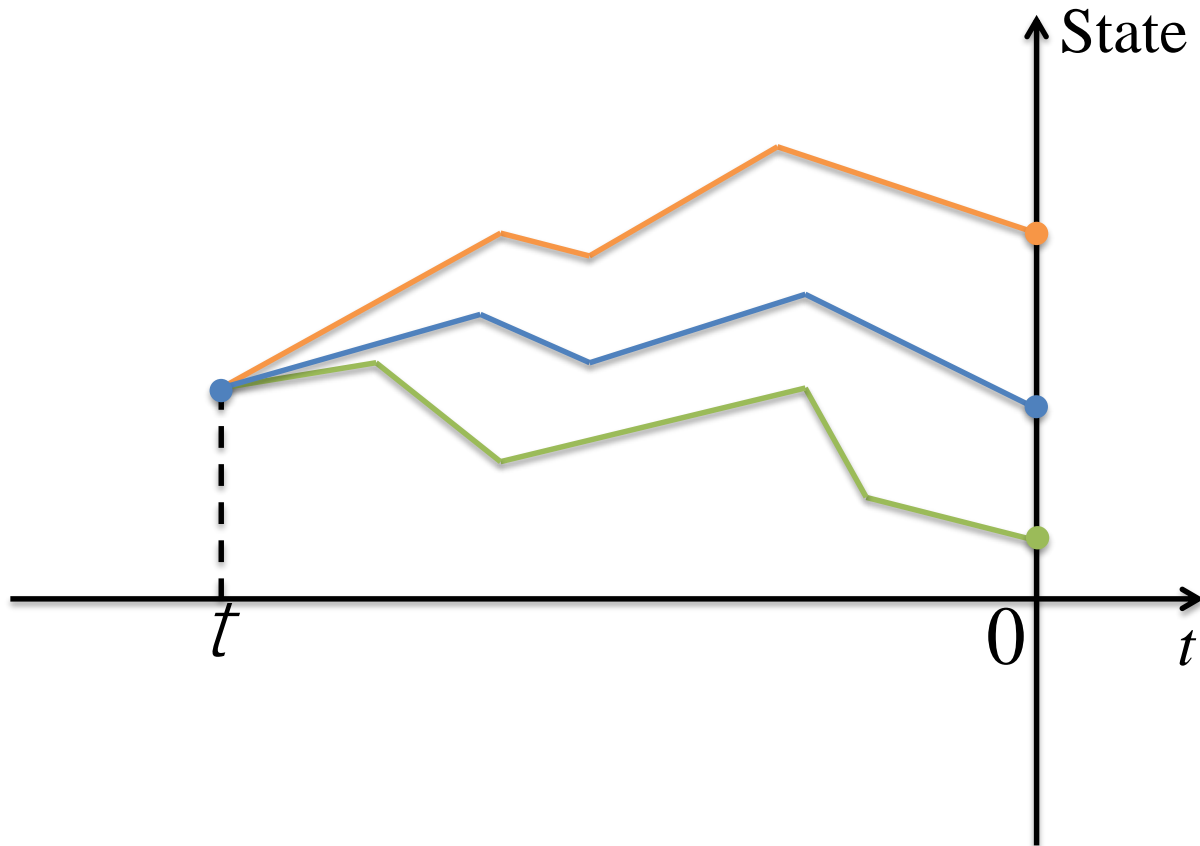
Dominated Coupling From The Past



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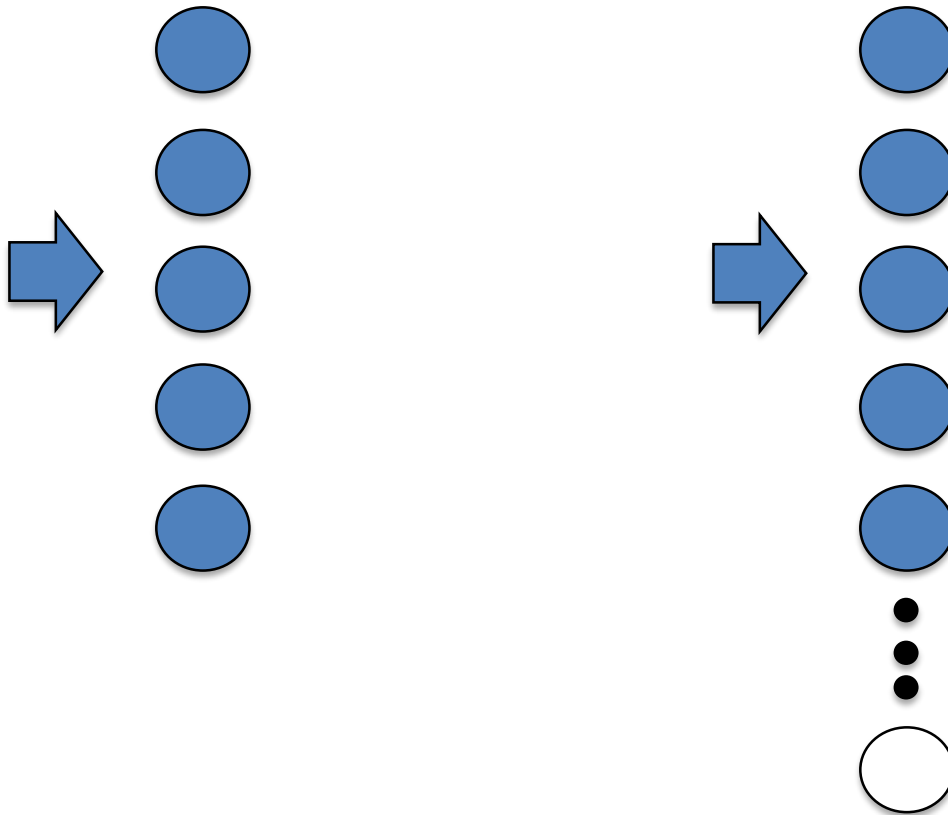


Dominated Coupling From The Past



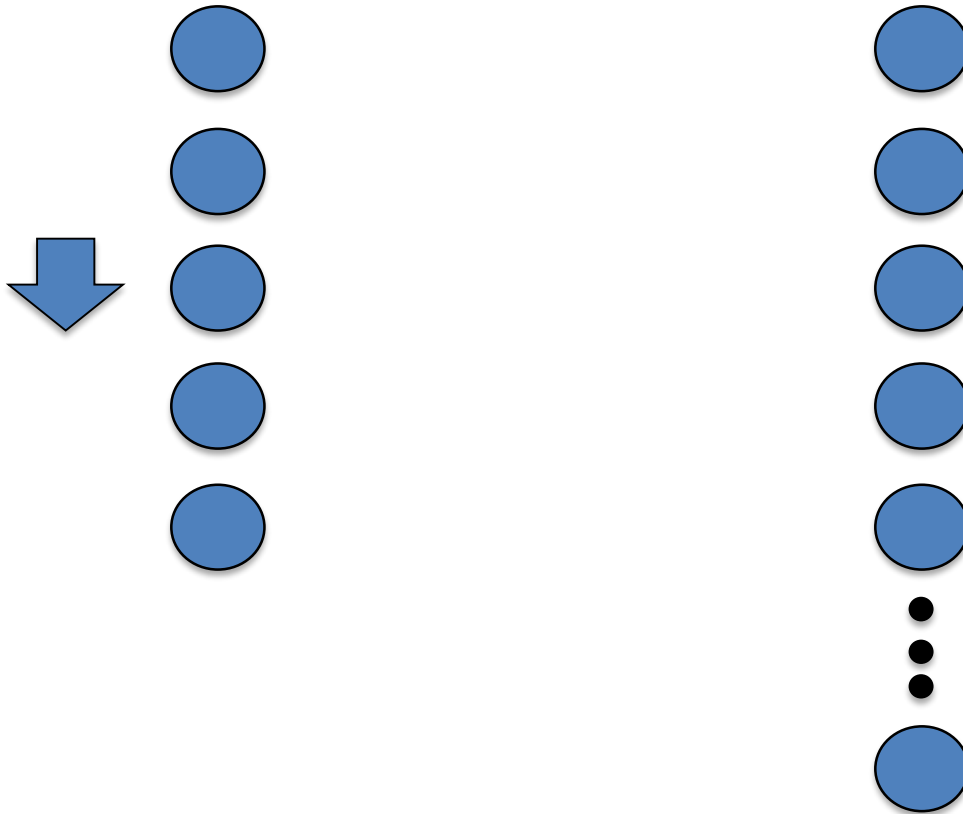
Loss queue: DCFTP

- Bounding process (upper and lower)
 - Upperbound: infinite server queue



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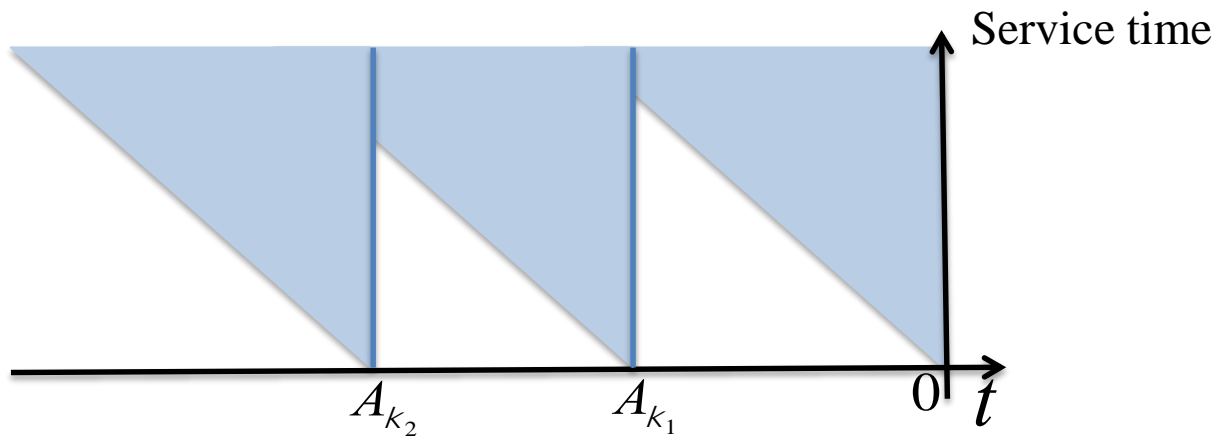


Loss queue: DCFTP

- Bounding process (upper and lower)
 - Upperbound: infinite server queue

$$Q^L(t, y) \leq Q^\forall(t, y) \quad \& \quad E^L(t) = E^\forall(t)$$

$$\supset W^L(t) \leq W^\forall(t)$$

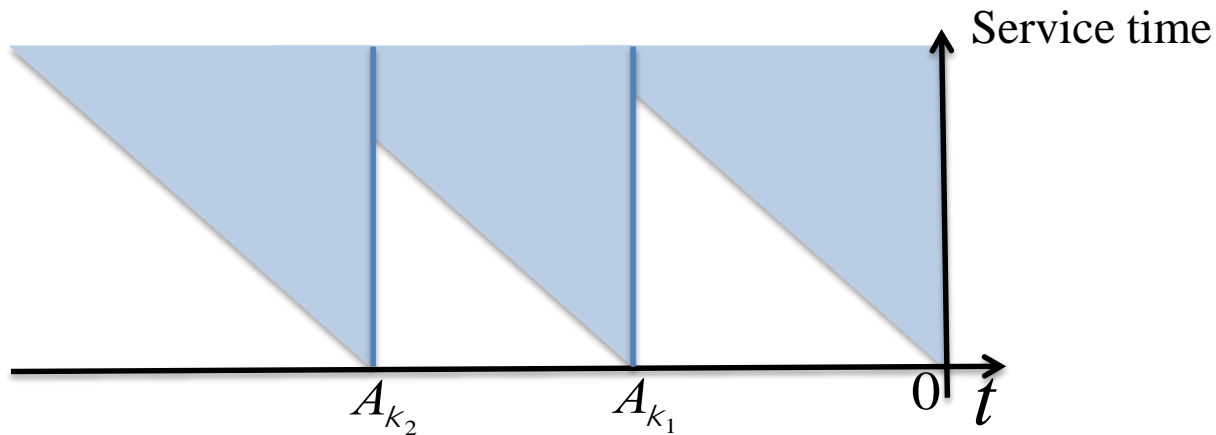


Loss queue: DCFTP

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$$\supset W^L(t) \leq W^\ddagger(t)$$



- Lower bound: 0? → not efficient!

Loss queue: DCFTP

- Coalescence time: $t + R^\ddagger(t)$
 where $R^\ddagger(t)$ largest remaining service time at time t

$$Q^\ddagger(t, 0) \in C, R^\ddagger(t) \in |t|$$

$$\max_{t \in t \in t + R^\ddagger(t)} Q^\ddagger(t, 0) \in C$$

